

Syntactic parsing.¹

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¹With materials used from "Speech and Language Processing", D. Jurafsky and J. H. Martin.

Constituents

- Morphology - words
- Syntax - in Greek *setting out together or arrangement*.
- Semantics - meaning

Constituents

- Constituents: groups of words behaving as a single units
 - Examples of NP constituents:
 - Harry the Horse
 - a high-class spot
 - such as Mindy's
 - the Broadway coppers
 - the reason he comes into the Hot Box
 - they
 - three parties from Brooklyn

Constituents

- Constituents appear in similar syntactic contexts
 - **three parties from Brooklyn** arrive. . .
 - a high-class spot **such as Mindy's** attracts. . .
 - **the Broadway coppers** love. . .
 - **they** sit
- Constituents can be moved inside the sentence:
 - **On September seventeenth**, I'd like to fly from Atlanta to Denver
 - I'd like to fly **on September seventeenth** from Atlanta to Denver
 - I'd like to fly from Atlanta to Denver **on September seventeenth**

Examples of grammar

NP → *Det Nominal*

NP → *ProperNoun*

Nominal → *Noun* | *Nominal Noun*

Det → *a*

Det → *the*

Noun → *flight*

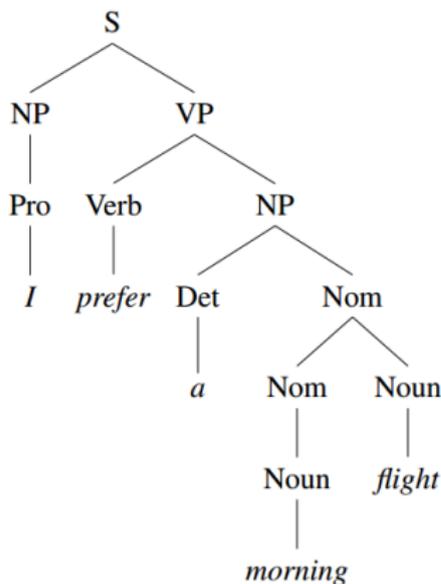
- Rules, generating language, may be nested.
- Symbols:
 - terminal: express final words
 - a, the, flight, etc.
 - non-terminal (express language abstractions)
 - NP, VP, PP, etc.
- Non-terminals, generating terminals, correspond to parts of speech.

More example rules

- $S \rightarrow NP VP$ (I prefer a morning flight)
- $VP \rightarrow Verb NP$ (prefer a morning flight)
- $VP \rightarrow Verb NP PP$ (leave Boston in the morning)
- $VP \rightarrow Verb PP$ (leaving on Thursday)
- $PP \rightarrow Preposition NP$ (from Los Angeles)

Parse trees

- Sample parse tree:



- Bracketed notation of parse trees (common in datasets)
 - [S [NP [Pro I]] [VP [V **prefer**] [NP [Det **a**] [Nom [N **morning**] [Nom [N **flight**]]]]]

Definitions

N a set of **non-terminal symbols** (or **variables**)
 Σ a set of **terminal symbols** (disjoint from N)
 R a set of **rules** or productions, each of the form $A \rightarrow \beta$,
 where A is a non-terminal,
 β is a string of symbols from the infinite set of strings $(\Sigma \cup N)^*$
 S a designated **start symbol** and a member of N

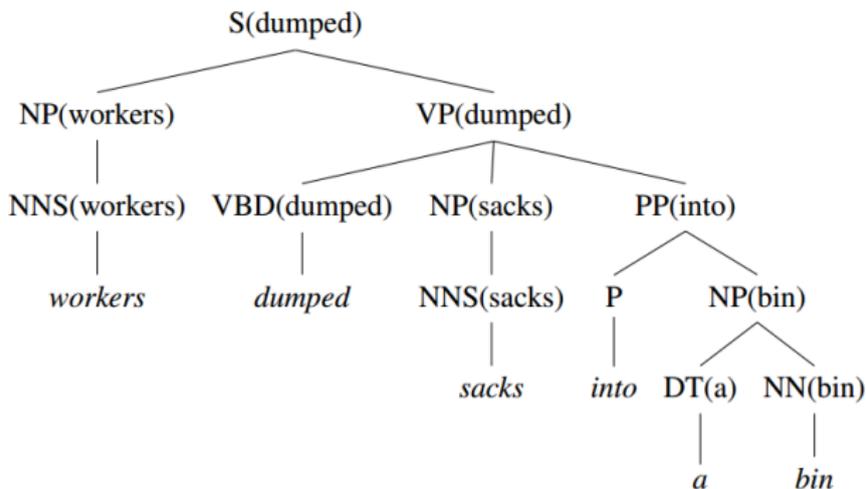
Notation:

- Capital letters like A , B , and S Non-terminals
- S The start symbol
- ε - empty symbol
- Lower-case Greek letters like α , β , and γ strings drawn from $(\Sigma \cup N)^*$
- Lower-case Roman letters like u , v , and w strings of terminals

Applications of grammar

- Language generation
 - A language, generated by grammar G is a subset of strings from Σ^* that can be derived using rules of G .
- Parsing of existing sentences
 - Syntactic parsing - mapping from a string of words to its parse tree
- Treebank - syntactically annotated corpus.
 - Penn Treebank project has treebanks from the Brown, Switchboard, ATIS, and Wall Street Journal corpora of English
 - We can extract grammar from treebank
 - it will be flat (many long rules), so post-processing is needed

Syntax parsing with head words



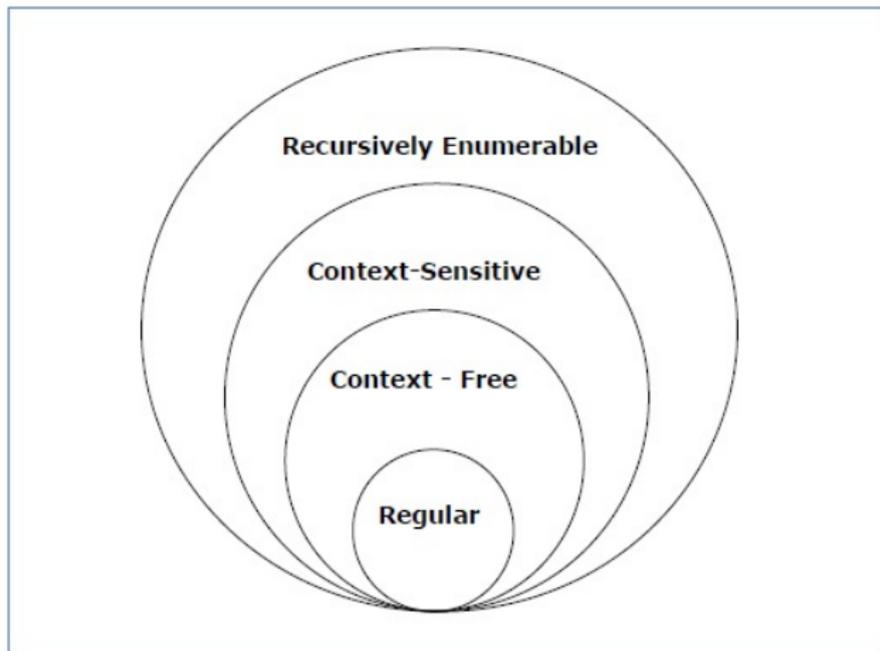
- Can incorporate head word into grammar
- More commonly: make regular parse and then assign head words with simple rules.

Grammars equivalence

- Grammars are weakly equivalent, if they generate the same set of strings.
- Grammars are strongly equivalent, if they generate the same set of strings and every sentence has the same parse tree
 - up to renaming non-terminals

Chomsky grammar classification

Chomsky grammar classification



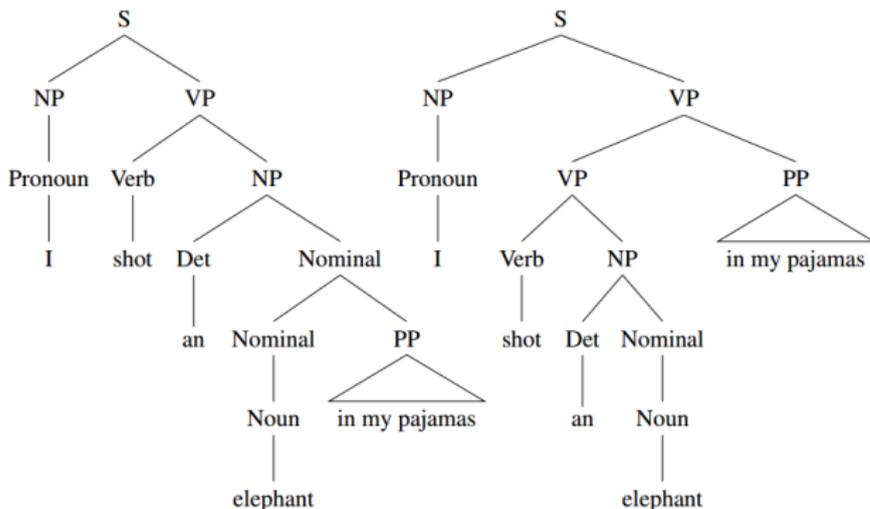
Chomsky grammar classification

- Regular language:
 - allowed rules: $X \rightarrow a$ or $X \rightarrow aY$, $X, Y \in N$, $a \in T$
 - $S \rightarrow \varepsilon$ is allowed only if S doesn't appear on the right side of any rule
 - recognizable by finite state automation or regular expressions.
- Context-free grammars:
 - allowed rules: $X \rightarrow \gamma$, $X \in N$, $\gamma \in (T \cup N)^*$
- Context-sensitive grammars:
 - $\alpha B \beta \rightarrow \alpha \gamma \beta$, where $B \in N$, $\alpha, \gamma, \beta \in (T \cup N)^*$
 - α, β may be empty, γ must be non-empty
 - $S \rightarrow \varepsilon$ is allowed only if S doesn't appear on the right side of any rule
- recursively enumerable grammar - no restrictions on rules
 - generate the languages that are recognized by a Turing machine.

Ambiguity

Examples of structural ambiguity:

- We saw [the Eiffel Tower flying to Paris].
- We saw [the Eiffel Tower] flying to Paris.
- [old [men and women]] <-> [old men] and [women]



Chomsky normal form

- Every context free grammar can be produced in Chomsky normal form (CNF)
- CYK (Cocke-Kasami-Younger) parser deals only with CNF context free grammars.
- In CNF all rules can be only:
 - $A \rightarrow v$ (non-terminal to terminal)
 - $A \rightarrow BC$ (non-terminal to 2 non-terminals)
- Example of CNF conversion:
 - $A \rightarrow B C D$ may be converted to
 - $A \rightarrow B X$
 - $X \rightarrow C D$

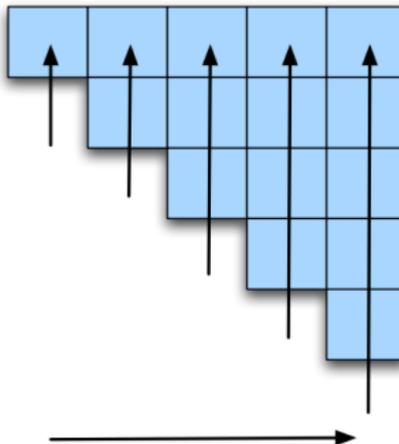
Conversion to CNF

- Mix of terminals and non-terminals on right hand side:
 - $INF\text{-}VP \rightarrow to\ VP$ replace by
 - $INF\text{-}VP \rightarrow TO\ VP$
 - $TO \rightarrow to$
- Single non-terminal on right hand side $A \rightarrow B$:
 - if A can be converted to B by some sequence of rules
 - for every rule $A \rightarrow \gamma$ add $B \rightarrow \gamma$
- More than 2 non-terminals on right hand side $A \rightarrow BC\gamma$
 - replace with:
 - $A \rightarrow X\gamma$
 - $X \rightarrow BC$

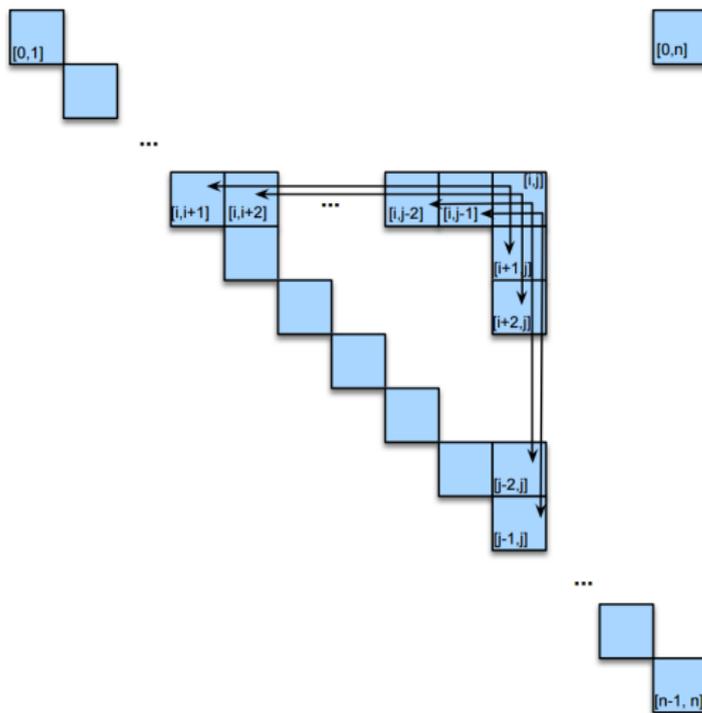
CYK matrix

- Positions of the sentence:
 - 0 Book 1 that 2 flight 3.
- Work with $(N + 1) \times (N + 1)$ matrix M
 - upper-triangular position are needed
 - $M[i, j]$ stands for parsing of sentence $[i:j]$
 - $M[i, j]$ if formed from $M[i, k]$ and $M[k, j]$ recurrently

	Book	the	flight	through	Houston
S, VP, Verb Nominal, Noun [0,1]	[0,2]	[0,3]	[0,4]	[0,5]	
	Det	NP		NP	
	[1,2]	[1,3]	[1,4]	[1,5]	
		Nominal, Noun		Nominal	
		[2,3]	[2,4]	[2,5]	
			Prep	PP	
			[3,4]	[3,5]	
				NP, Proper- Noun	
				[4,5]	



CYK illustration



CYK algorithm

function CKY-PARSE(*words*, *grammar*) **returns** *table*

for $j \leftarrow$ **from** 1 **to** LENGTH(*words*) **do**

for all $\{A \mid A \rightarrow \text{words}[j] \in \text{grammar}\}$
 $\text{table}[j-1, j] \leftarrow \text{table}[j-1, j] \cup A$

for $i \leftarrow$ **from** $j-2$ **downto** 0 **do**

for $k \leftarrow i+1$ **to** $j-1$ **do**

for all $\{A \mid A \rightarrow BC \in \text{grammar} \text{ and } B \in \text{table}[i, k] \text{ and } C \in \text{table}[k, j]\}$
 $\text{table}[i, j] \leftarrow \text{table}[i, j] \cup A$

- This is a recognizer
- To recognize sentence it needs to find S in matrix $[0, N]$
- To make it a parser we need to add backtracking
 - add from which elements each new element can be derived

Restoring original grammar

- We can augment CYK parser to deal with rules $A \rightarrow B$
- Rules $A \rightarrow BC$, $A \rightarrow wC$, $A \rightarrow Cw$ can be restored by merging CNF rules

Partial parsing

- Partial (shallow) parsing - parse only lower layers of tree
 - chunk sentence into segments
 - [NP The morning flight] [PP from] [NP Denver] [VP has arrived.]
 - some words may not be covered
 - e.g. when only NP are interesting: [NP The morning flight] from [NP Denver] has arrived.
- Applications: information extraction from most informative sentence parts.
- Chunking is performed with sequence labelling (aka POS tagging)
- IOB notation:
 - B-beginning of next chunk
 - I-inside current chunk
 - O-outside any chunk

Examples of partial parsing

- [NP The morning flight] [PP from] [NP Denver] [VP has arrived.]

B	I	I	B	B	B	I
The	morning	flight	from	Denver	has	arrived

- [NP The morning flight] from [NP Denver] has arrived.

B	I	I	O	B	O	O
The	morning	flight	from	Denver	has	arrived

- We can automatically deduce the endings of each chunk

Examples of partial parsing with chunk type

- [NP The morning flight] [PP from] [NP Denver] [VP has arrived.]

B_NP	I_NP	I_NP	B_PP	B_NP	B_VP	I_VP
The	morning	flight	from	Denver	has	arrived

- [NP The morning flight] from [NP Denver] has arrived.

B_NP	I_NP	I_NP	O	B_NP	O	O
The	morning	flight	from	Denver	has	arrived

- We can automatically deduce the endings of each chunk
- Total number of tags $2n + 1$, where n is the number of chunk types.

Training

- Use treebanks to generate training set
- Use sequence labelling algorithms
- Features:
 - look on words at positions -2,-1,0,1,2
 - get
 - words themselves
 - parts of speech
 - previous classifications at positions -2,-1.

Evaluation

$$\text{Precision} = \frac{\text{Number of correct chunks given by system}}{\text{Total number of chunks given by system}}$$

$$\text{Recall} = \frac{\text{Number of correct chunks given by system}}{\text{Total number of actual chunks in the text}}$$

Overall measure (weighted harmonic mean)

$$F_{\beta} = \frac{1}{\frac{\beta^2}{\beta^2+1} \frac{1}{R} + \frac{1}{\beta^2+1} \frac{1}{P}}$$

or simply (uniform harmonic mean)

$$F_1 = \frac{1}{\frac{1}{2} \frac{1}{R} + \frac{1}{2} \frac{1}{P}}$$

State-of-the-art results

- State-of-the-art results:
 - NP chunking, English language: 0.96
 - NP,VP,PP,ADVP,SBAR,ADJP chunking, English language: 0.92 - 0.94
- Problems:
 - POS tagging accuracy
 - inconsistencies of correct labels generation from parse trees
 - ambiguities:
 - [NP Late arrivals and departures] are commonplace during winter.
 - [NP Late arrivals] and [NP cancellations] are commonplace during winter.
 - semantic information is needed to make correct chunking!

PCFG

Probabilistic context free grammar:

N	set of non-terminal symbols
Σ	set of terminal symbols
R	set of rules $A \rightarrow \beta [p]$, $A \in N$, $\beta \in (\Sigma \cup N)^*$
S	start symbol

- We can estimate probabilities from treebanks.
- Independence assumption: rules application probabilities don't depend on already applied rules.

Probabilistic parsing

- S -observed sentence, T -unobserved parse tree
- We are interested to find most probable parse:

$$T = \arg \max_T p(T|S)$$

- Probability of sentence:

$$P(S) = \sum_T p(S|T)$$

Probabilistic CYK algorithm

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function PROBABILISTIC-CKY(words, grammar) returns most probable parse
                                                    and its probability
  for j ← from 1 to LENGTH(words) do
    for all A;  $A \rightarrow \text{words}[j] \in \text{grammar}$ 
       $\text{table}[j-1, j, A] \leftarrow P(A \rightarrow \text{words}[j])$ 
    for i ← from j - 2 downto 0 do
      for k ← i + 1 to j - 1 do
        for all A;  $A \rightarrow BC \in \text{grammar}$ ,
          and  $\text{table}[i, k, B] > 0$  and  $\text{table}[k, j, C] > 0$ 
          if ( $\text{table}[i, j, A] < P(A \rightarrow BC) \text{table}[i, k, B] \text{table}[k, j, C]$ ) then
             $\text{table}[i, j, A] \leftarrow P(A \rightarrow BC) \text{table}[i, k, B] \text{table}[k, j, C]$ 
             $\text{back}[i, j, A] \leftarrow k, B, C$ 
  return BUILD_TREE( $\text{back}[1, \text{LENGTH}(\text{words}), S]$ ),  $\text{table}[1, \text{LENGTH}(\text{words}), S]$ 

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